

Rioplattense 2019/L2/3

Evan Chen

TWITCH SOLVES ISL

Episode 132

Problem

Let n be a positive integer. Let S be a multiset of $2n + 1$ numbers less or equal than 2^n with the following property: the product of any n numbers of S divides the product of the $n + 1$ remaining numbers of S . Prove that S has at least 2 equal numbers.

Video

<https://youtu.be/3X6YcQybomM>

External Link

<https://aops.com/community/p28630932>

Solution

We start by analyzing a particular prime:

Claim. Let p be a fixed prime, and let

$$e_0 \leq \cdots \leq e_{2n}$$

be the ν_p 's of the elements of S . There exists an index k such that

$$e_{k+1} = e_{k+2} \cdots = e_{k+2n-e_0}.$$

Proof. Write the above inequality as

$$\sum_{i=1}^n (e_n - e_i) + \sum_{i=n+1}^{2n} (e_i - e_n) \leq e_0$$

where

$$\begin{aligned} 0 &= e_n - e_n \leq e_n - e_{n-1} \leq e_n - e_{n-2} \leq \cdots \leq e_n - e_1 \\ 0 &\leq e_{n+1} - e_n \leq e_{n+2} - e_n \leq \cdots \leq e_{2n} - e_n. \end{aligned}$$

If we set $A = \sum_{i=1}^n (e_n - e_i)$ and $B = \sum_{i=n+1}^{2n} (e_i - e_n)$, then in fact the smallest $n - A$ summands of A and the smallest $n - B$ summands of B must all be zero. These produce the $A + B \leq 2n - e_0$ required indices. \square

We now proceed in two cases, by induction on n . The base cases $n = 0, 1, 2$ are clear, so assume $n \geq 3$.

- Assume there is a prime p that doesn't divide every element of S (i.e. $e_0 = 0$ for some prime p in the claim's notation). Then by the above, in fact there is one element of S not divisible by p , say x , and all other elements of S have the same exponent of p . Pick any $y \in S$ with $y \neq x$ and apply induction hypothesis onto $\frac{1}{p} \cdot (S \setminus \{x, y\})$.
- Otherwise, assume we're in a situation where the primes p_1, \dots, p_m appear with nonzero minimum multiplicities u_1, \dots, u_m , respectively. Obviously we may assume $\sum u_i \leq n - 1$, otherwise even the GCD of all the elements is at least $2^{u_1 + \dots + u_m}$. In particular, $\max(u_i) \leq n - 1$; also (as $u_i \geq 1$) we have $m \leq n - 1$.

Imagine the following algorithm. Write all elements of S on the blackboard. Then for $i = 1, \dots, m$, take the index promised by the claim for the prime p_i , and draw an X on all $u_i + 1$ elements not corresponding to those e_i 's. If, after the algorithm ends, there are at least two numbers not X'ed out, we win.

However, the total number of X's is $\sum (u_i + 1) \leq (n - 1) + m = 2n - 2$.