

# China 2021/6

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TWITCH SOLVES ISL

Episode 123

## Problem

Solve over positive integers the functional equation

$$f(f(x) + y) \mid x + f(y).$$

## Video

<https://youtu.be/orvTuytr6uk>

## External Link

<https://aops.com/community/p19116098>

## Solution

There are three families of solutions:

- $f$  is the identity function;
- $f(n) = 1$  for  $n \geq 2$ , where  $f(1)$  is any positive integer.
- $f(n) = 2$  for  $n$  even,  $f(n) = 1$  for odd  $n \geq 3$ , and  $f(1)$  is any odd positive integer.

The verification is easy, so we prove these are the only solution. The proof is split into two main cases.

**Case where  $f$  is not injective** Let

$$T = \min_{f(a)=f(b), a < b} (b - a).$$

Then all outputs of  $f$  are eventually divisors of  $T$ , as

$$f(y + f(a)) = f(y + f(b)) \mid \gcd(f(y) + a, f(y) + b) \mid b - a \quad \forall y > 0.$$

**Claim.** We have  $T \leq 2$ .

*Proof.* If  $T > 2$ , then look any  $T$  consecutive outputs. They were supposed to be distinct divisors of  $T$ , impossible for  $T > 2$ .  $\square$

If  $T = 1$ , then for any  $n \geq 2$ , let  $y = n - 1$  and let  $x$  be a large integer. Then we have  $f(n) \mid x + f(n - 1)$  for all large  $x$ , forcing  $f(n) = 1$ .

If  $T = 2$ , then it follows  $f$  must alternate between 1 and 2 eventually. Again,  $f(n) \leq 2$  follows for  $n \geq 2$  as in the previous case, by letting  $y = n - 1$  and  $x$  be large with  $f(x) = 1$ . We can then quickly verify that only the situation where  $f(\text{even}) = \text{even}$  bears solutions, the ones we claimed earlier.

**Case where  $f$  is injective** Let  $c = f(1)$ . Taking  $x = 1$  then gives

$$f(y + c) \leq f(y) + 1.$$

This implies

$$f(n) \leq \frac{1}{c}n + \max\{f(1), \dots, f(c)\}.$$

For  $f$  to be injective, we must then have  $c = 1$ . Finally,  $f(y + 1) \leq f(y) + 1$  together with  $f$  injective then forces  $f$  to be the identity function, by induction.