

# USA TST 2024 Solutions

## United States of America — Team Selection Test

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## §0 Problems

1. Find the smallest constant  $C > 1$  such that the following statement holds: for every integer  $n \geq 2$  and sequence of non-integer positive real numbers  $a_1, a_2, \dots, a_n$  satisfying

$$\frac{1}{a_1} + \frac{1}{a_2} + \dots + \frac{1}{a_n} = 1,$$

it's possible to choose positive integers  $b_i$  such that

- (i) for each  $i = 1, 2, \dots, n$ , either  $b_i = \lfloor a_i \rfloor$  or  $b_i = \lfloor a_i \rfloor + 1$ ; and  
 (ii) we have

$$1 < \frac{1}{b_1} + \frac{1}{b_2} + \dots + \frac{1}{b_n} \leq C.$$

2. Let  $ABC$  be a triangle with incenter  $I$ . Let segment  $AI$  intersect the incircle of triangle  $ABC$  at point  $D$ . Suppose that line  $BD$  is perpendicular to line  $AC$ . Let  $P$  be a point such that  $\angle BPA = \angle PAI = 90^\circ$ . Point  $Q$  lies on segment  $BD$  such that the circumcircle of triangle  $ABQ$  is tangent to line  $BI$ . Point  $X$  lies on line  $PQ$  such that  $\angle IAX = \angle XAC$ . Prove that  $\angle AXP = 45^\circ$ .
3. Let  $n > k \geq 1$  be integers and let  $p$  be a prime dividing  $\binom{n}{k}$ . Prove that the  $k$ -element subsets of  $\{1, \dots, n\}$  can be split into  $p$  classes of equal size, such that any two subsets with the same sum of elements belong to the same class.
4. Find all integers  $n \geq 2$  for which there exists a sequence of  $2n$  pairwise distinct points  $(P_1, \dots, P_n, Q_1, \dots, Q_n)$  in the plane satisfying the following four conditions:
- (i) no three of the  $2n$  points are collinear;
  - (ii)  $P_i P_{i+1} \geq 1$  for all  $i = 1, 2, \dots, n$ , where  $P_{n+1} = P_1$ ;
  - (iii)  $Q_i Q_{i+1} \geq 1$  for all  $i = 1, 2, \dots, n$ , where  $Q_{n+1} = Q_1$ ; and
  - (iv)  $P_i Q_j \leq 1$  for all  $i = 1, 2, \dots, n$  and  $j = 1, 2, \dots, n$ .
5. Suppose  $a_1 < a_2 < \dots < a_{2024}$  is an arithmetic sequence of positive integers, and  $b_1 < b_2 < \dots < b_{2024}$  is a geometric sequence of positive integers. Find the maximum possible number of integers that could appear in both sequences, over all possible choices of the two sequences.
6. Solve over  $\mathbb{R}$  the functional equation

$$f(xf(y)) + f(y) = f(x + y) + f(xy).$$

## §1 Solutions to Day 1

### §1.1 USA TST 2024/1, proposed by Merlijn Staps

Available online at <https://aops.com/community/p29409075>.

#### Problem statement

Find the smallest constant  $C > 1$  such that the following statement holds: for every integer  $n \geq 2$  and sequence of non-integer positive real numbers  $a_1, a_2, \dots, a_n$  satisfying

$$\frac{1}{a_1} + \frac{1}{a_2} + \dots + \frac{1}{a_n} = 1,$$

it's possible to choose positive integers  $b_i$  such that

(i) for each  $i = 1, 2, \dots, n$ , either  $b_i = \lfloor a_i \rfloor$  or  $b_i = \lfloor a_i \rfloor + 1$ ; and

(ii) we have

$$1 < \frac{1}{b_1} + \frac{1}{b_2} + \dots + \frac{1}{b_n} \leq C.$$

¶ **Answer.** The answer is  $C = \frac{3}{2}$ .

¶ **Lower bound.** Note that if  $a_1 = \frac{4n-3}{2n-1}$  and  $a_i = \frac{4n-3}{2}$  for  $i > 1$ , then we must have  $b_1 \in \{1, 2\}$  and  $b_i \in \{2n-2, 2n-1\}$  for  $i > 1$ . If we take  $b_1 = 2$  then we obtain

$$\frac{1}{b_1} + \frac{1}{b_2} + \dots + \frac{1}{b_n} \leq \frac{1}{2} + (n-1) \cdot \frac{1}{2n-2} = 1,$$

whereas if we take  $b_1 = 1$  then we obtain

$$\frac{1}{b_1} + \frac{1}{b_2} + \dots + \frac{1}{b_n} \geq 1 + (n-1) \cdot \frac{1}{2n-1} = \frac{3n-2}{2n-1}.$$

This shows that  $C \geq \frac{3n-2}{2n-1}$ , and as  $n \rightarrow \infty$  this shows that  $C \geq \frac{3}{2}$ .

¶ **Upper bound.** For  $0 \leq k \leq n$ , define

$$c_i = \sum_{j=1}^k \frac{1}{\lfloor a_j \rfloor} + \sum_{j=k+1}^n \frac{1}{\lfloor a_j \rfloor + 1}.$$

Note that  $c_0 < c_1 < \dots < c_n$  and

$$c_0 < \frac{1}{a_1} + \frac{1}{a_2} + \dots + \frac{1}{a_n} = 1 < c_n.$$

This means there exists a unique value of  $k$  for which  $c_{k-1} < 1 < c_k$ . For this  $k$  we have

$$1 < c_k = c_{k-1} + \frac{1}{(\lfloor a_k \rfloor)(\lfloor a_k \rfloor + 1)} < 1 + \frac{1}{1 \cdot 2} = \frac{3}{2}.$$

Therefore we may choose  $b_i = \lfloor a_i \rfloor$  for  $i \leq k$  and  $b_i = \lfloor a_i \rfloor + 1$  for  $i > k$ .

**Remark.** The solution can be phrased in the following “motion-based” way. Imagine starting with all floors (corresponding to  $c_0$ ), then changing each floor to a ceiling one by one until after  $n$  steps every floor is a ceiling (arriving at  $c_n$ ). As we saw,  $c_0 < 1 < c_n$ , but  $c_0 < \dots < c_n$ . Moreover, each discrete step increases the sum by at most

$$\frac{1}{\lfloor a_i \rfloor} - \frac{1}{\lfloor a_i \rfloor + 1} \leq \frac{1}{2}$$

and so the changing sum must be in the interval  $[1, 3/2]$  at some point.

¶ **Upper bound (alternate).** First suppose  $a_i < 2$  for some  $i$ . Assume without loss of generality  $i = 1$  here. Let  $b_1 = 1$  and  $b_i = \lfloor a_i \rfloor + 1$  for all other  $i$ . Then

$$\begin{aligned} 1 &< \frac{1}{b_1} + \dots + \frac{1}{b_n} = 1 + \frac{1}{\lfloor a_2 \rfloor + 1} + \dots + \frac{1}{\lfloor a_n \rfloor + 1} \\ &< \left( \frac{1}{2} + \frac{1}{a_1} \right) + \frac{1}{a_2} + \dots + \frac{1}{a_n} = \frac{3}{2}. \end{aligned}$$

Now suppose  $a_i > 2$  always. Then  $\frac{a_i}{\lfloor a_i \rfloor} < \frac{3}{2}$ , so

$$1 = \frac{1}{a_1} + \dots + \frac{1}{a_n} < \frac{1}{\lfloor a_1 \rfloor} + \dots + \frac{1}{\lfloor a_n \rfloor} < \frac{3}{2} \left( \frac{1}{a_1} + \dots + \frac{1}{a_n} \right) = \frac{3}{2}.$$

Therefore we may let  $b_i = \lfloor a_i \rfloor$  for all  $i$ .

**Remark.** The original proposal asked to find the optimal  $C$  for a fixed  $n$ . The answer is  $\frac{3n-2}{2n-1}$ , i.e. the lower bound construction in the solution is optimal.

## §1.2 USA TST 2024/2, proposed by Luke Robitaille

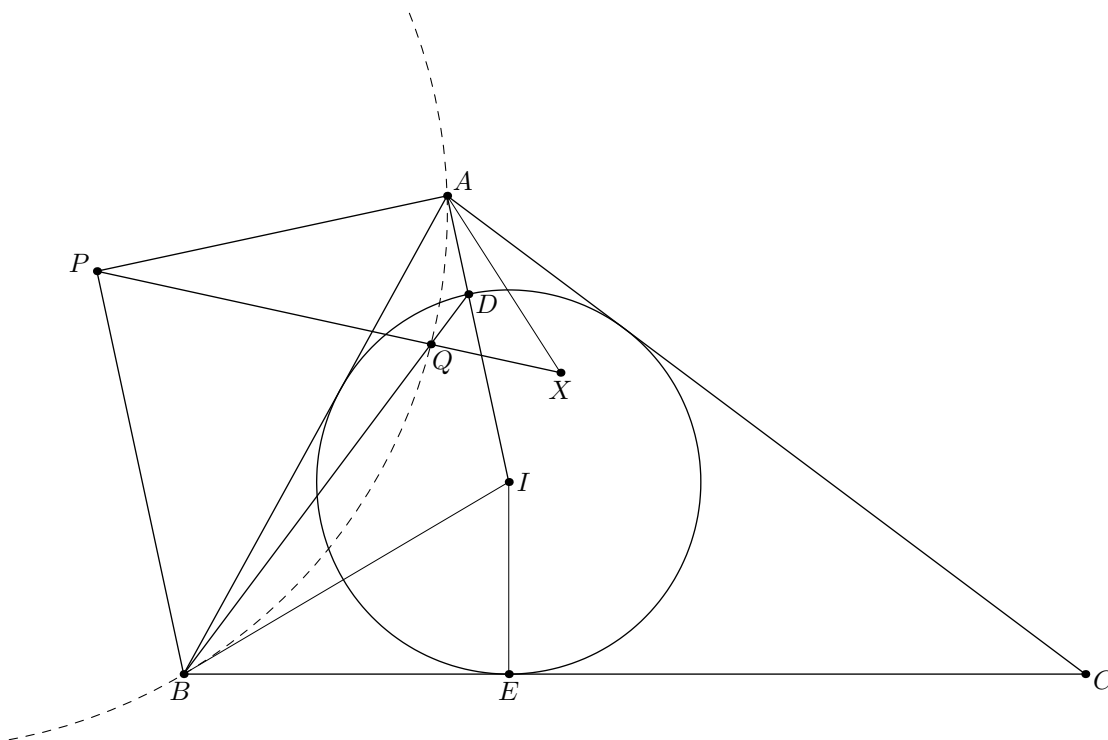
Available online at <https://aops.com/community/p29409083>.

### Problem statement

Let  $ABC$  be a triangle with incenter  $I$ . Let segment  $AI$  intersect the incircle of triangle  $ABC$  at point  $D$ . Suppose that line  $BD$  is perpendicular to line  $AC$ . Let  $P$  be a point such that  $\angle BPA = \angle PAI = 90^\circ$ . Point  $Q$  lies on segment  $BD$  such that the circumcircle of triangle  $ABQ$  is tangent to line  $BI$ . Point  $X$  lies on line  $PQ$  such that  $\angle IAX = \angle XAC$ . Prove that  $\angle AXP = 45^\circ$ .

We show several approaches.

### ¶ First solution, by author.



**Claim** — We have  $BP = BQ$ .

*Proof.* For readability, we split the proof into three unconditional parts.

- We translate the condition  $\overline{BD} \perp \overline{AC}$ . It gives  $\angle DBA = 90^\circ - A$ , so that

$$\begin{aligned} \angle DBI &= \left| \frac{B}{2} - (90^\circ - A) \right| = \frac{|A - C|}{2} \\ \angle BDI &= \angle DBA + \angle BAD = (90^\circ - A) + \frac{A}{2} = 90^\circ - \frac{A}{2}. \end{aligned}$$

Hence, letting  $r$  denote the inradius, we can translate  $\overline{BD} \perp \overline{AC}$  into the following trig condition:

$$\sin \frac{B}{2} = \frac{r}{BI} = \frac{DI}{BI} = \frac{\sin \angle DBI}{\sin \angle BDI} = \frac{\sin \frac{|A-C|}{2}}{\sin \left(90^\circ - \frac{A}{2}\right)}.$$

- The length of  $BP$  is given from right triangle  $APB$  as

$$BP = BA \cdot \sin \angle PAB = BA \cdot \sin \left(90^\circ - \frac{A}{2}\right).$$

- The length of  $BQ$  is given from the law of sines on triangle  $ABQ$ . The tangency gives  $\angle BAQ = \angle DBI$  and  $\angle BQA = 180^\circ - \angle ABI = 180^\circ - \angle IBE$  and thus

$$BQ = BA \cdot \frac{\sin \angle BAQ}{\sin \angle AQB} = BA \cdot \frac{\sin \angle DBI}{\sin \angle ABI} = BA \cdot \frac{\sin \frac{|A-C|}{2}}{\sin \frac{B}{2}}.$$

The first bullet implies the expressions in the second and third bullet for  $BP$  and  $BQ$  are equal, as needed.  $\square$

**Remark.** In the above proof, one does not actually need to compute  $\angle DBI = \frac{|A-C|}{2}$ . The proof works equally leaving that expression intact as  $\sin \angle DBI$  in place of  $\sin \frac{|A-C|}{2}$ .

Now we can finish by angle chasing. We have

$$\angle PBQ = \angle PBA + \angle ABD = \frac{A}{2} + 90^\circ - A = 90^\circ - \frac{A}{2}.$$

Then

$$\angle BPQ = \frac{180^\circ - \angle PBQ}{2} = 45^\circ + \frac{A}{4},$$

so  $\angle APQ = 90^\circ - \angle BPQ = 45^\circ - \frac{A}{4}$ . Also, if we let  $J$  be the incenter of  $IAC$ , then  $\angle PAJ = 90^\circ + \frac{A}{4}$ , and clearly  $X$  lies on line  $AJ$ . Then  $\angle APQ + \angle PAJ = 135^\circ < 180^\circ$ , so  $X$  lies on the same side of  $AP$  as  $Q$  and  $J$  (by the parallel postulate). Therefore  $\angle AXP = 180^\circ - 135^\circ = 45^\circ$ , as desired.

**Remark.** The problem was basically written backwards by starting from the  $BD \perp AC$  condition, turning that into trig, and then contriving  $P$  and  $Q$  so that the  $BD \perp AC$  condition implied  $BP = BQ$ .

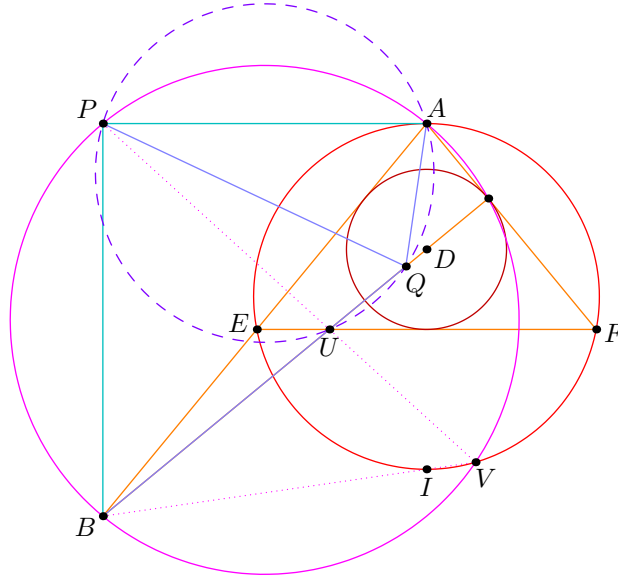
¶ **Second solution, by Jeffrey Kwan.** We prove the following restatement:

Consider isosceles triangle  $AEF$  with  $AE = AF$  and incenter  $D$ . Let  $B$  be the point on ray  $AE$  such that  $BD \perp AF$ , and let  $P$  be the projection of  $B$  onto the line through  $A$  parallel to  $EF$ . Let  $I$  be the point diametrically opposite  $A$  in the circumcircle of  $AEF$ , and let  $Q$  be the point on line  $BD$  such that  $BI$  is tangent to the circumcircle of  $AQB$ . Then  $\angle APQ = 45^\circ - \angle A/4$ .

First note that  $\angle DFE = 45^\circ - \angle A/4$ , so it suffices to show that  $\overline{PQ} \parallel \overline{DF}$ . Let  $U = \overline{BD} \cap \overline{EF}$ , and let  $V = BI \cap (AEF)$ . Observe that:

- $P$  and  $V$  both lie on the circle with diameter  $AB$ , so  $\angle BVP = \angle PAB = 90^\circ - \angle A/2$ .
- We have  $\angle EVB = \angle EVI = \angle A/2 = \angle DUF = \angle BUE$ . Hence  $BEUV$  is cyclic.

Now  $\angle BVU = \angle AEU = 90^\circ - \angle A/2 = \angle BVP$ , so  $\overline{PUV}$  are collinear.



From the tangency condition, we have that  $\angle AQB = 180^\circ - \angle ABI$ , which implies that

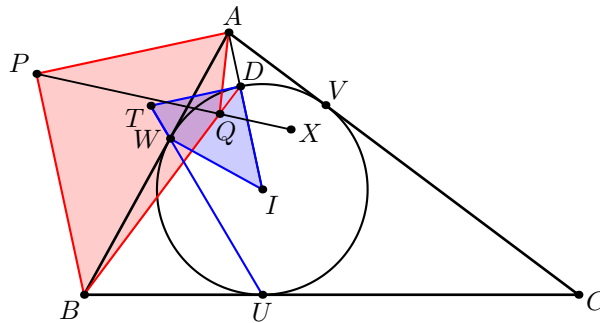
$$\angle AQU + \angle APU = \angle AQB + \angle APV = (180^\circ - \angle ABI) + \angle ABI = 180^\circ,$$

and so  $APUQ$  is cyclic. Finally, note that  $D$  is the orthocenter of  $\triangle AUF$ , which implies that

$$\angle APQ = \angle AUQ = \angle AUD = \angle AFD = \angle DFE.$$

This forces  $\overline{PQ} \parallel \overline{DF}$ , as desired.

¶ **Third solution by Pitchayut Saengrungskongka and Maxim Li.** We provide yet another proof that  $BP = BQ$ .



Let the incircle be  $\omega$  and touch  $BC$  and  $AB$  at point  $U$  and  $W$ . Let the tangent to  $\omega$  at  $D$  meet  $UW$  at  $T$ . Notice that  $T$  is the pole of  $BD$  with respect to  $\omega$ , so  $IT \perp BD$ . Now, we make the following critical claim, which in particular implies  $BP = BQ$ .

**Claim** — Quadrilaterals  $DIWT$  and  $PBQA$  are inversely similar.

*Proof.* This follows from four angle relations.

- $\angle IDT = \angle BPA = 90^\circ$ .
- $\angle TIW = \angle ABQ$ .
- $\angle DIT = \angle IAC = \angle BAI = \angle ABP$ .
- $\angle ITW = \angle QBI = \angle QAB$ . □

With  $BP = BQ$  obtained, one finishes with the same angle chasing as in the first solution.



### §1.3 USA TST 2024/3, proposed by Ankan Bhattacharya

Available online at <https://aops.com/community/p29409068>.

#### Problem statement

Let  $n > k \geq 1$  be integers and let  $p$  be a prime dividing  $\binom{n}{k}$ . Prove that the  $k$ -element subsets of  $\{1, \dots, n\}$  can be split into  $p$  classes of equal size, such that any two subsets with the same sum of elements belong to the same class.

Let  $\sigma(S)$  denote the sum of the elements of  $S$ , so that

$$P(x) := \sum_{\substack{S \subseteq \{1, \dots, n\} \\ |S|=k}} x^{\sigma(S)}$$

is the generating function for the sums of  $k$ -element subsets of  $\{1, \dots, n\}$ .

By Legendre's formula,

$$\nu_p \left( \binom{n}{k} \right) = \sum_{r=1}^{\infty} \left( \left\lfloor \frac{n}{p^r} \right\rfloor - \left\lfloor \frac{k}{p^r} \right\rfloor - \left\lfloor \frac{n-k}{p^r} \right\rfloor \right)$$

so there exists a positive integer  $r$  with

$$\left\lfloor \frac{n}{p^r} \right\rfloor - \left\lfloor \frac{k}{p^r} \right\rfloor - \left\lfloor \frac{n-k}{p^r} \right\rfloor > 0.$$

The main claim is the following:

**Claim** —  $P(x)$  is divisible by

$$\Phi_{p^r}(x) = x^{(p-1)p^{r-1}} + \dots + x^{p^{r-1}} + 1.$$

Before proving this claim, we will show how it solves the problem. It implies that there exists a polynomial  $Q$  with integer coefficients satisfying

$$\begin{aligned} P(x) &= \Phi_{p^r}(x)Q(x) \\ &= (x^{(p-1)p^{r-1}} + \dots + x^{p^{r-1}} + 1)Q(x). \end{aligned}$$

Let  $c_0, c_1, \dots$  denote the coefficients of  $P$ , and define

$$s_i = \sum_{j \equiv i \pmod{p^r}} c_j.$$

Then it's easy to see that

$$\begin{aligned} s_0 &= s_{p^{r-1}} = \dots = s_{(p-1)p^{r-1}} \\ s_1 &= s_{p^{r-1}+1} = \dots = s_{(p-1)p^{r-1}+1} \\ &\vdots \\ s_{p^{r-1}-1} &= s_{2p^{r-1}-1} = \dots = s_{p^r-1}. \end{aligned}$$

This means we can construct the  $p$  classes by placing a set with sum  $z$  in class  $\left\lfloor \frac{z \pmod{p^r}}{p^{r-1}} \right\rfloor$ .

Now we present two ways to prove the claim.

¶ **First proof of claim.** There's a natural bijection between  $k$ -element subsets of  $\{1, \dots, n\}$  and binary strings consisting of  $k$  zeroes and  $\ell$  ones: the set  $\{a_1, \dots, a_k\}$  corresponds to the string which has zeroes at positions  $a_1, \dots, a_k$ . Moreover, the inversion count of this string is simply  $(a_1 + \dots + a_k) - \frac{1}{2}k(k+1)$ , so we only deal with these inversion counts (equivalently, we are factoring  $x^{\frac{k(k+1)}{2}}$  out of  $P$ ).

Recall that the generating function for these inversion counts is given by the  $q$ -binomial coefficient

$$P(x) = \frac{(x-1) \cdots (x^{k+\ell} - 1)}{[(x-1) \cdots (x^k - 1)] \times [(x-1) \cdots (x^\ell - 1)]}.$$

By choice of  $r$ , the numerator of  $P(x)$  has more factors of  $\Phi_{p^r}(x)$  than the denominator, so  $\Phi_{p^r}(x)$  divides  $P(x)$ .

**Remark.** Here is a proof that  $P(x)$  is divisible by  $\Phi_{p^r}(x)$  for some  $r$  using the  $q$ -binomial formula, without explicitly identifying  $r$ . We know that  $P(x)$  is the product of several cyclotomic polynomials, and that  $P(1)$  is a multiple of  $p$ . Thus there is a factor  $\Phi_q(x)$  for which  $\Phi_q(1)$  is a multiple of  $p$ , which is equivalent to  $q$  being a power of  $p$ .

¶ **Second proof of claim.** Note that  $P(x)$  is the coefficient of  $y^k$  in the polynomial

$$Q(x, y) := (1 + xy)(1 + x^2y) \cdots (1 + x^ny).$$

Let  $a$  be the remainder when  $n$  is divided by  $p^r$ , and let  $b$  be the remainder when  $k$  is divided by  $p^r$ ; then we have  $a < b$  by the choice of  $r$ . Let  $q = \lfloor n/p^r \rfloor$  so  $n = p^r q + a$ . Consider taking  $x$  to be a primitive  $p^r$ th root of unity, say  $\omega$ . Then

$$Q(\omega, y) = [(1 + \omega y)(1 + \omega^2 y) \cdots (1 + \omega^{p^r} y)]^q (1 + \omega y)(1 + \omega^2 y) \cdots (1 + \omega^a y).$$

Now  $\omega, \omega^2, \dots, \omega^{p^r}$  are all the  $p^r$ th roots of unity, each exactly once; then we can see that

$$\begin{aligned} & (1 + \omega y)(1 + \omega^2 y) \cdots (1 + \omega^{p^r} y) \\ &= (1 - \omega(-y))(1 - \omega^2(-y)) \cdots (1 - \omega^{p^r}(-y)) \\ &= 1 - (-y)^{p^r}, \end{aligned}$$

so

$$Q(\omega, y) = (1 - (-y)^{p^r})^q (1 + \omega y)(1 + \omega^2 y) \cdots (1 + \omega^a y).$$

In particular, for any  $m$ , if the coefficient of  $y^m$  in  $Q(\omega, x)$  is nonzero, then  $m$  must be congruent to one of  $0, 1, \dots, a \pmod{p^r}$ . Therefore the coefficient of  $y^k$  in  $Q(\omega, y)$  is zero. This means that  $P(\omega) = 0$  whenever  $\omega$  is a primitive  $p^r$ th root of unity, which proves the claim.

## §2 Solutions to Day 2

### §2.1 USA TST 2024/4, proposed by Ray Li

Available online at <https://aops.com/community/p29643081>.

#### Problem statement

Find all integers  $n \geq 2$  for which there exists a sequence of  $2n$  pairwise distinct points  $(P_1, \dots, P_n, Q_1, \dots, Q_n)$  in the plane satisfying the following four conditions:

- (i) no three of the  $2n$  points are collinear;
- (ii)  $P_i P_{i+1} \geq 1$  for all  $i = 1, 2, \dots, n$ , where  $P_{n+1} = P_1$ ;
- (iii)  $Q_i Q_{i+1} \geq 1$  for all  $i = 1, 2, \dots, n$ , where  $Q_{n+1} = Q_1$ ; and
- (iv)  $P_i Q_j \leq 1$  for all  $i = 1, 2, \dots, n$  and  $j = 1, 2, \dots, n$ .

¶ **Answer.** Even integers only.

¶ **Proof that even  $n$  work.** If we ignore the conditions that the points are pairwise distinct and form no collinear triples, we may take

$$P_{2i+1} = (0.51, 0), \quad P_{2i} = (-0.51, 0), \quad Q_{2i+1} = (0, 0.51), \quad Q_{2i} = (0, -0.51).$$

The distances  $P_i P_{i+1}$  and  $Q_i Q_{i+1}$  are  $1.02 > 1$ , while the distances  $P_i Q_j$  are  $0.51\sqrt{2} < 1$ . We may then perturb each point by a small amount to ensure that the distance inequalities still hold and have the points in general position.

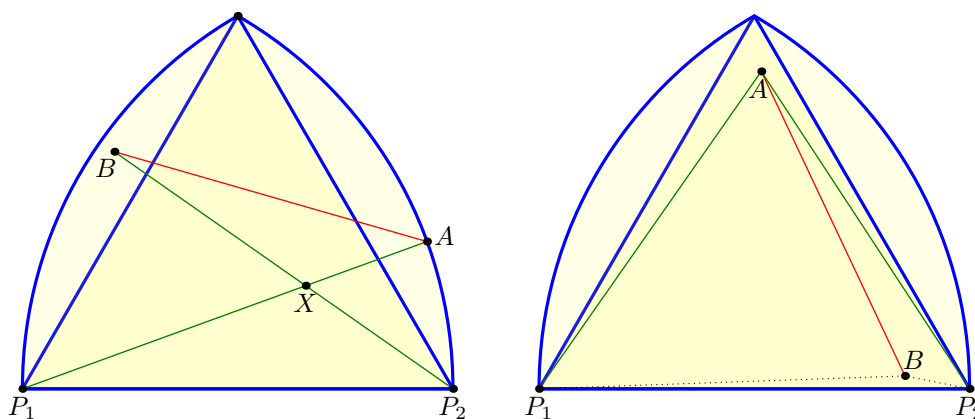
¶ **Proof that odd  $n$  do not work.** The main claim is the following.

**Claim** — For  $1 \leq i \leq n$ , points  $Q_i$  and  $Q_{i+1}$  must lie on opposite sides of line  $P_1 P_2$ .

To isolate the geometry component of the problem, we rewrite the claim in the following contrapositive form, without referencing the points  $Q_i$  and  $Q_{i+1}$ :

#### Lemma

Suppose  $A$  and  $B$  are two points such that  $\max(P_1 A, P_1 B, P_2 A, P_2 B) \leq 1$ , and moreover  $A$  and  $B$  lie on the same side of line  $P_1 P_2$ . Further assume no three of  $\{P_1, P_2, A, B\}$  are collinear. Then  $AB < 1$ .



*Proof of lemma.* Suppose for the sake of contradiction that  $A$  and  $B$  lie on the same side of  $P_1P_2$ . The convex hull of these four points is either a quadrilateral or a triangle.

- If the convex hull is a quadrilateral, assume WLOG that the vertices are  $P_1P_2AB$  in order. Let  $X$  denote the intersection of segments  $P_1A$  and  $P_2B$ . Then

$$1 + AB = P_1P_2 + AB < P_1X + XP_2 + AX + XB = P_1A + P_2B \leq 2.$$

- Otherwise, assume WLOG that  $B$  is in the interior of triangle  $P_1P_2A$ . Since  $\angle P_1BA + \angle P_2BA = 360^\circ - \angle P_1BP_2 > 180^\circ$ , at least one of  $\angle P_1BA$  and  $\angle P_2BA$  is obtuse. Assume WLOG the former angle is obtuse; then  $AB < P_1A \leq 1$ .  $\square$

**Remark.** Another proof of the lemma can be found by replacing segment  $AB$  with the intersection of this line on the boundary of the blue region above, which does not decrease the distance. In other words, one can assume WLOG that  $A$  and  $B$  lie on either segment  $AB$  or one of the two circular arcs. One then proves that  $AB \leq 1$ , and that for equality to occur, one of  $A$  and  $B$  must lie on segment  $P_1P_2$ . However, this approach seems to involve a fair bit more calculation.

Yet another clever approach uses the trivia-fact that a **Reuleaux triangle** happens to have constant width.

In any case, it's important to realize that this claim is *not trivial*; while it looks like it is easy to prove, it is not, owing to the two near-equality cases.

It follows from the claim that  $Q_i$  is on the same side of line  $P_1P_2$  as  $Q_1$  if  $i$  is odd, and on the opposite side if  $i$  is even. Since  $Q_1 = Q_{n+1}$ , this means the construction is not possible when  $n$  is odd.

**Remark.** The fact that  $n$  cannot be odd follows from Theorem 3 of **EPTAS for Max Clique on Disks and Unit Balls**. In the language of that paper, if  $G$  is a unit ball graph, then the *induced odd cycle parking number* of  $\bar{G}$  is at most 1.

In earlier versions of the proposed problem, the points were not necessarily distinct to make the even  $n$  case nicer, but this resulted in annoying boundary conditions for the odd  $n$  case.

## §2.2 USA TST 2024/5, proposed by Ray Li

Available online at <https://aops.com/community/p29642892>.

### Problem statement

Suppose  $a_1 < a_2 < \dots < a_{2024}$  is an arithmetic sequence of positive integers, and  $b_1 < b_2 < \dots < b_{2024}$  is a geometric sequence of positive integers. Find the maximum possible number of integers that could appear in both sequences, over all possible choices of the two sequences.

¶ **Answer.** 11 terms.

¶ **Construction.** Let  $a_i = i$  and  $b_i = 2^{i-1}$ .

¶ **Bound.** We show a  $\nu_p$ -based approach communicated by Derek Liu, which seems to be the shortest one. At first, we completely ignore the geometric sequence  $b_i$  and focus only on the arithmetic sequence.

**Claim** — Let  $p$  be any prime, and consider the sequence

$$\nu_p(a_1), \nu_p(a_2), \dots, \nu_p(a_{2024}).$$

Set  $C := \lfloor \log_p(2023) \rfloor$ . Then there are at most  $C + 2$  different values in this sequence.

*Proof.* By scaling, assume the  $a_i$  do not have a common factor, so that  $a_i = a + di$  where  $\gcd(a, d) = 1$ .

- If  $p \mid d$ , then  $p \nmid a$  and  $\nu_p(a_i)$  is constant.
- Otherwise, assume  $p \nmid d$ . We will in fact prove that every term in the sequence is contained in  $\{0, 1, \dots, C\}$  with at most one exception.

Define  $M := \max_i \nu_p(a_i)$ . If  $M \leq C$ , there is nothing to prove. Otherwise, fix some index  $m$  such that  $\nu_p(a_m) = M$ . We know  $\nu_p(i - m) \leq C$  since  $|i - m| \leq 2023$ . But now for any other index  $i \neq m$ ;

$$\begin{aligned} \nu_p(d(i - m)) &= \nu_p(i - m) \leq C < M = \nu_p(a_m) \\ \implies \nu_p(a_i) &= \nu_p(a_m + d(i - m)) = \nu_p(i - m) \leq C. \end{aligned}$$

so  $\nu_p(a_m)$  is the unique exceptional term of the sequence exceeding  $C$ . □

**Remark.** The bound in the claim is best possible by taking  $a_i = p^M + (i - 1)$  for any  $M > C$ . Then indeed, the sequence  $\nu_p(a_i)$  takes on values in  $\{0, 1, \dots, C\}$  for  $i > 1$  while  $\nu_p(a_1) = M$ .

Back to the original problem with  $(b_i)$ . Consider the common ratio  $r \in \mathbb{Q}$  of the geometric sequence  $(b_i)$ . If  $p$  is any prime with  $\nu_p(r) \neq 0$ , then every term of  $(b_i)$  has a different  $\nu_p$ . So there are two cases to think about.

- Suppose any  $p \geq 3$  has  $\nu_p(r) \neq 0$ . Then there are at most

$$2 + \log_p(2023) = 2 + \log_3(2023) \approx 8.929 < 11$$

overlapping terms, as needed.

- Otherwise, suppose  $r$  is a power of 2 (in particular,  $r \geq 2$  is an integer). We already have an upper bound of 12; we need to improve it to 11.

As in the proof before, we may assume WLOG by scaling down by a power of 2 that the common difference of  $a_i$  is odd. (This may cause some  $b_i$  to become rational non-integers, but that's OK. However, unless  $\nu_2(a_i)$  is constant, the  $a_i$  will still be integers.)

Then in order for the bound of 12 to be achieved, the sequence  $\nu_2(a_i)$  must be contained in  $\{0, 1, \dots, 10, M\}$  for some  $M \geq 11$ . In particular, we only need to work with  $r = 2$ .

Denote by  $b$  the unique odd-integer term in the geometric sequence, which must appear among  $(a_i)$ . Then  $2b$  appears too, so the common difference of  $a_i$  is at most  $b$ .

But if  $(a_i)$  is an arithmetic progression of integers that contains  $b$  and has common difference at most  $b$ , then no term of the sequence can ever exceed  $b + 2023 \cdot b = 2024b$ . Hence  $2^M b$  cannot appear for any  $M \geq 11$ . This completes the proof.

**Remark.** There are several other approaches to the problem, but most take some time to execute. The primary issue is that the common difference of the  $a_i$ 's could share prime factors with the common ratio in the  $b_i$ 's, which means that merely trying to write out a lot of modular arithmetic equations leads to a lot of potential technical traps that are not pleasant to defuse.

One unusual thing is that many solutions end up proving a bound of 12 (in the case the common ratio is 2) and then having to adjust it to 11 later.

### §2.3 USA TST 2024/6, proposed by Milan Haiman

Available online at <https://aops.com/community/p29642897>.

#### Problem statement

Solve over  $\mathbb{R}$  the functional equation

$$f(xf(y)) + f(y) = f(x + y) + f(xy).$$

In addition to all constant functions,  $f(x) \equiv x + 1$  clearly works too. We prove these are the only solutions. The solution that follows is by the original proposer.

Let  $P(x, y)$  denote the given assertion.

**Claim 2.1** — If  $f$  is periodic, then  $f$  is constant.

*Proof.* Let  $f$  have period  $d \neq 0$ . From  $P(x, y + d)$ , we have

$$f(x(y + d)) = f(x + y + d) - f(y + d) - f(xf(y + d)) = f(x + y) - f(y) - f(xf(y)).$$

Applying  $P(x, y)$  gives

$$f(x(y + d)) = f(xy).$$

In particular, taking  $y = 0$  yields that  $f(dx) = f(0)$ . Thus  $f$  is constant, as  $d \neq 0$ .  $\square$

**Claim 2.2** — For all real numbers  $x$  and  $y$ , we have  $f(f(x) + y) = f(f(y) + x)$ .

*Proof.* Applying  $P(f(x), y)$  and then  $P(y, x)$  gives us

$$\begin{aligned} f(f(x)f(y)) &= f(f(x) + y) + f(f(x)y) - f(y) \\ &= f(f(x) + y) + f(x + y) + f(xy) - f(x) - f(y). \end{aligned}$$

Now swapping  $x$  and  $y$  gives us  $f(f(x) + y) = f(f(y) + x)$ .  $\square$

**Claim 2.3** — If  $f$  is nonconstant, then  $f(f(x) + y) = f(x) + f(y)$  for all reals  $x, y$ .

*Proof.* Let  $x, y \in \mathbb{R}$  and let  $d := f(f(x) + y) - f(x) - f(y)$ . Let  $z$  be an arbitrary real number. By repeatedly applying [Claim 2.2](#), we have

$$\begin{aligned} f(z + f(f(x) + y)) &= f(f(z) + f(x) + y) \\ &= f(f(x) + f(z) + y) \\ &= f(x + f(f(z) + y)) \\ &= f(x + f(f(y) + z)) \\ &= f(f(x) + f(y) + z) \\ &= f(z + f(x) + f(y)). \end{aligned}$$

If  $d \neq 0$ , then  $f$  is periodic with period  $d$ , contradicting [Claim 2.1](#).  $\square$

**Claim 2.4** — If  $f$  is nonconstant, then  $f(0) = 1$  and  $f(x + 1) = f(x) + 1$ .

*Proof.* From  $P(z, 0)$  we have  $f(zf(0)) = f(z)$  for all real  $z$ . Then  $P(xf(0), y)$  and  $P(x, y)$  give us

$$\begin{aligned} f(xf(0) + y) &= f(y) + f(xf(0)f(y)) - f(xf(0)y) \\ &= f(y) + f(xf(y)) - f(xy) = f(x + y). \end{aligned}$$

If  $f(0) \neq 1$ , then  $xf(0) - x$  is a period of  $f$  for all  $x$ , violating [Claim 2.1](#). So we must have  $f(0) = 1$ .

Now putting  $x = 0$  in [Claim 2.3](#) gives  $f(x + 1) = f(x) + 1$ .  $\square$

**Claim 2.5** — If  $f$  is nonconstant, then  $f(x) + f(y) = f(x + y) + 1$ .

*Proof.* From  $P(x + 1, y)$  and [Claim 2.4](#), we have

$$f((x + 1)f(y)) = f(x + y + 1) + f(xy + y) - f(y) = f(x + y) + f(xy + y) - f(y) + 1.$$

Also, from [Claim 2.3](#) and  $P(x, y)$ , we have

$$f((x + 1)f(y)) = f(xf(y)) + f(y) = f(x + y) + f(xy).$$

Thus  $f(xy) = f(xy + y) - f(y) + 1$ . Replacing  $x$  with  $\frac{x}{y}$  gives the claim for all  $y \neq 0$  (whereas  $y = 0$  follows from [Claim 2.4](#)).  $\square$

**Claim 2.6** — If  $f$  is nonconstant, then  $f(x) \equiv x + 1$ .

*Proof.* We apply [Claim 2.3](#), [Claim 2.5](#), and [Claim 2.4](#):

$$f(f(x) + y) = f(x) + f(y) = f(x + y) + 1 = f(x + y + 1).$$

If  $f(x) \neq x + 1$  for some  $x$ , then [Claim 2.1](#) again gives a contradiction.  $\square$